

Graph Theory

Lecture 9: Kruskal's Algorithm

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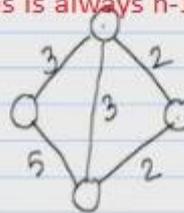
Minimum Spanning Trees

Spanning tree definition:

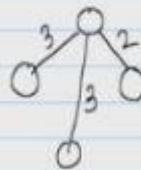
Minimum Spanning Trees (MST)

A spanning tree of a graph G is a subgraph that is a tree which includes all the vertices of G .

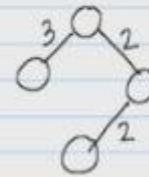
Note: no. of edges in a spanning tree of n vertices is always $n-1$



Weighted graph G



$$\text{Cost}(T_1) = 8$$

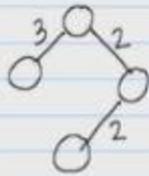


$$\text{Cost}(T_2) = 7$$

Minimum Spanning Tree definition

A minimum spanning tree (MST) of G is a spanning tree of G for which the sum of edge costs is minimum.

From previous slide spanning trees i.e. T_1 and T_2 T_2 is minimum weighted tree



$$\text{Cost}(T_2) = 7$$

T_2 is the unique minimum spanning tree of G

Theorem :

Algorithm : Kruskal's Algorithm for MST

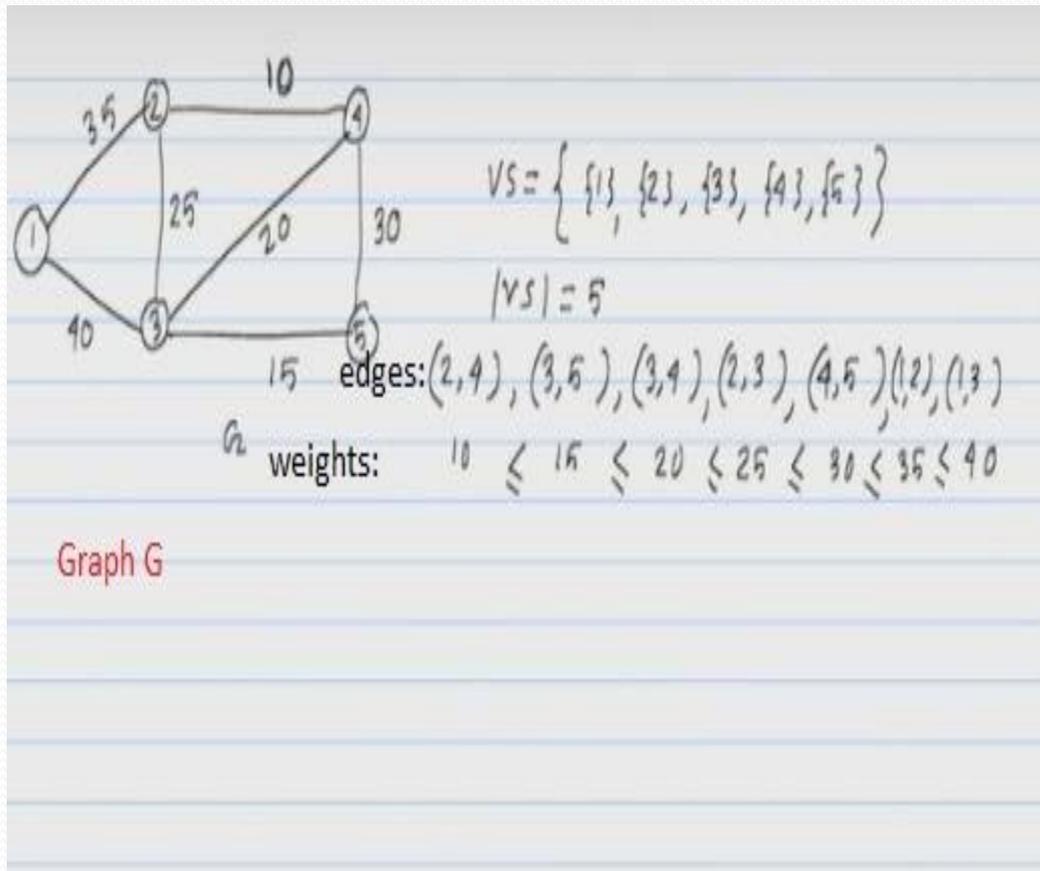
Kruskal's Algorithm input: $G = (V, E)$
output: A MST

$T = \emptyset$; As initially no node is spanned
 $VS = \emptyset$; VS = Vertex set for spanning tree

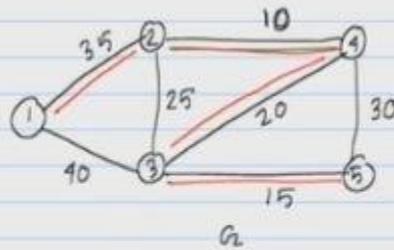
foreach $v \in V$, do add $\{v\}$ to VS;
order the edges (v, w) in non-decreasing order of weights and store in $Q = \text{queue}$

while $(|VS| > 1)$ { $|VS| = \text{cardinality or no. of vertices of the graph}$
 choose an edge (v, w) of lowest cost
 delete (v, w) from Q
 if v and w are in different sets w_1 & w_2 in VS
 { replace w_1 & w_2 in VS by $w_1 \cup w_2$;
 $T = T \cup (v, w)$;
 }
 }
} return T .

Example: For finding MST Using Kruskal's algorithm



Example(Continues...)



w1 and w2 are components or partitions

$$VS = \left\{ \overset{w_1}{\{1\}}, \overset{w_2}{\{2\}}, \{3\}, \{4\}, \{5\} \right\}$$

$$|VS| = 5$$

$$(\cancel{2,4}), (\cancel{3,5}), (\cancel{3,4}), (\cancel{2,3}), (\cancel{4,5}), (\cancel{1,3})$$

$$10 \leq 15 \leq 20 \leq 25 \leq 30 \leq 35 \leq 40$$

edges	action	VS
(2,4)	add	$\left\{ \{1\}, \overset{w_1}{\{2,4\}}, \overset{w_2}{\{3\}}, \{5\} \right\}$
(3,5)	add	$\left\{ \{1\}, \overset{w_2}{\{2,4\}}, \overset{w_1}{\{3,5\}} \right\}$
(3,4)	add	$\left\{ \overset{w_1}{\{1\}}, \overset{w_2}{\{2,3,4,5\}} \right\}$
(2,3)	reject	
(4,5)	reject	
(1,2)	add	$\left\{ \{1,2,3,4,5\} \right\}$

$T = \left\{ (2,4), (3,5), (3,4), (1,2) \right\}$ } Then T is MST, as is shown above in graph with red ^{supp.} $|VS| = 1$

Another Interpretation of Kruskal's algorithm

1. Treat n nodes of the graph as n independent trees of the forest
2. Join first the 2 vertices with minimum possible weights.
3. repeat step 2 until no vertices remains for traversal.

